

Superviz Plenary Meeting

WP4-TH4.1

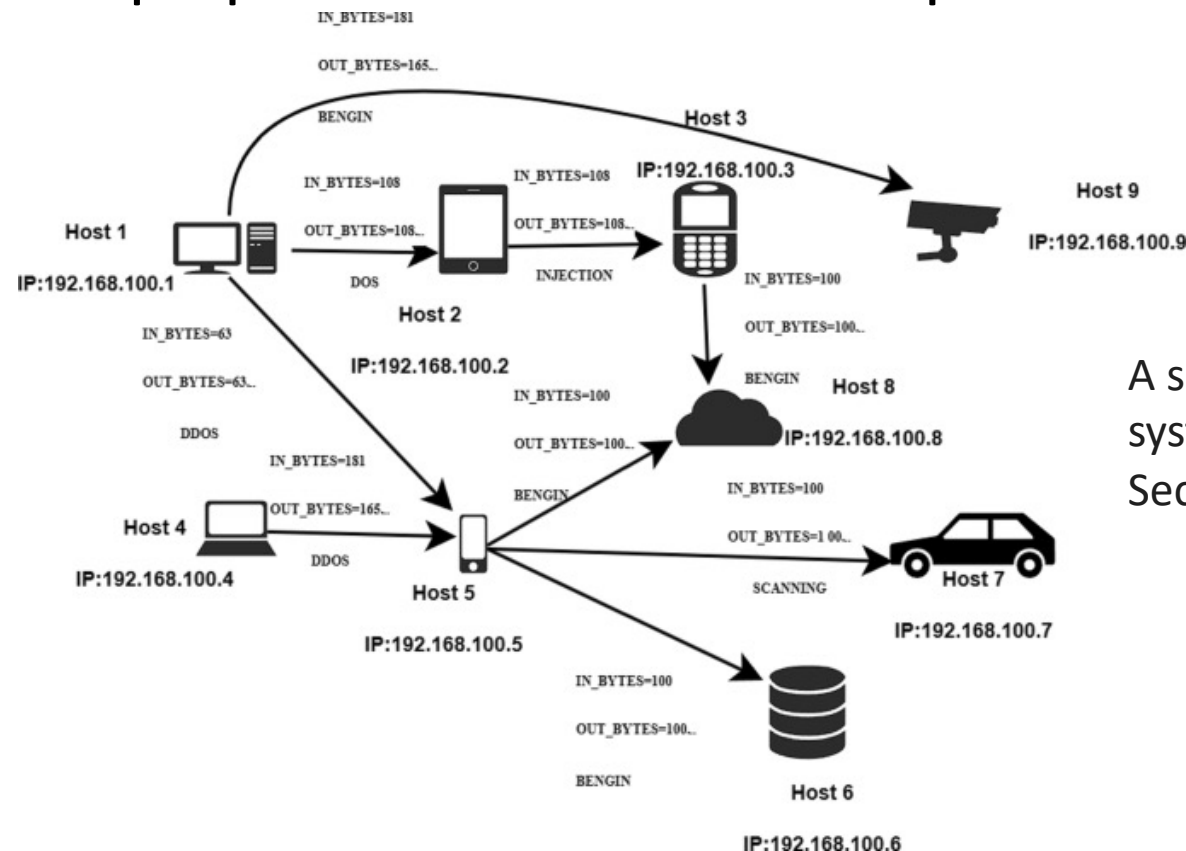
Yufei Han@INRIA PIRAT

March 11, 2025

@Campus Cyber, Paris

Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

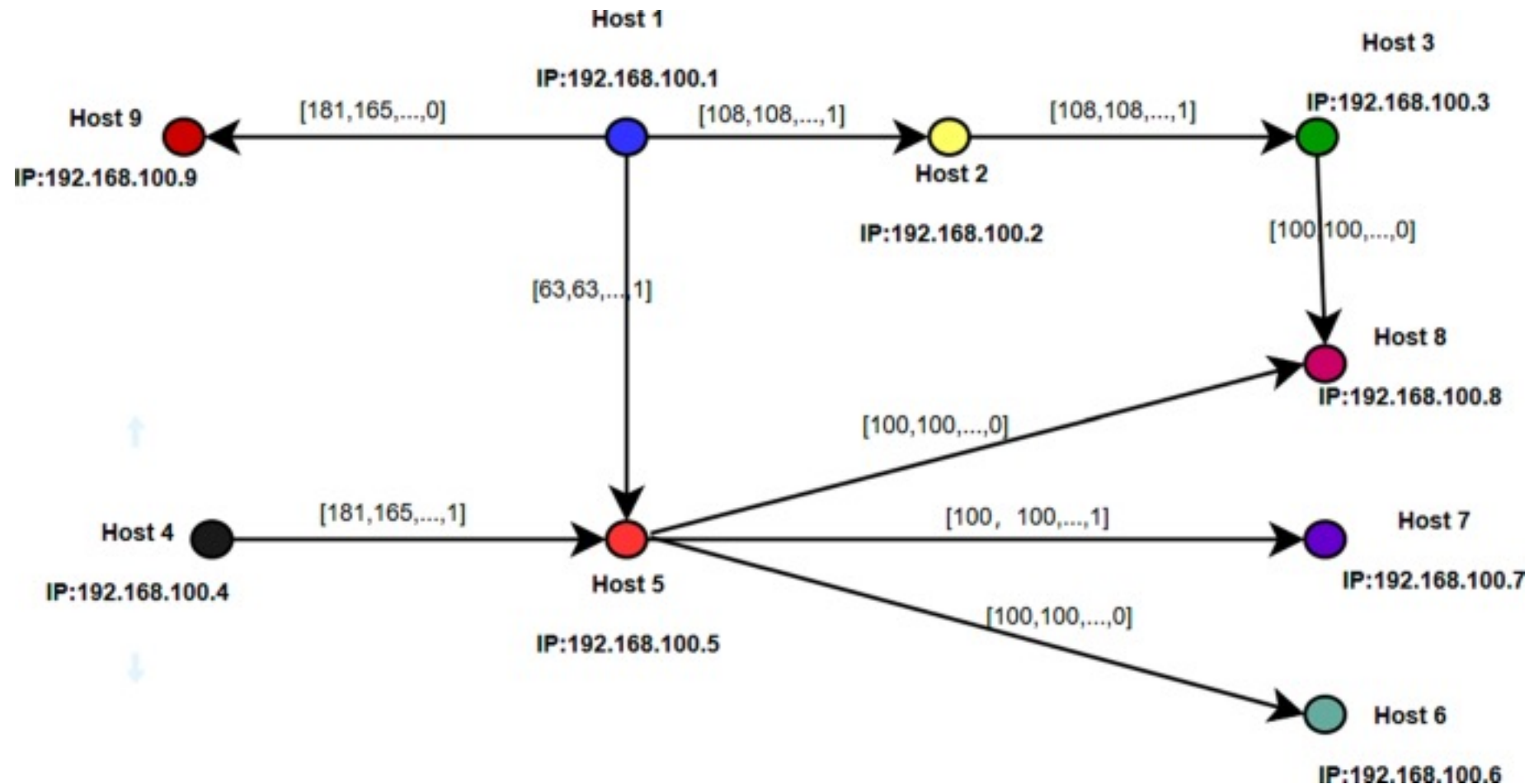
- Why Graph-based Intrusion Detection ?
- Graph provides a structural representation of cyber attack behaviour



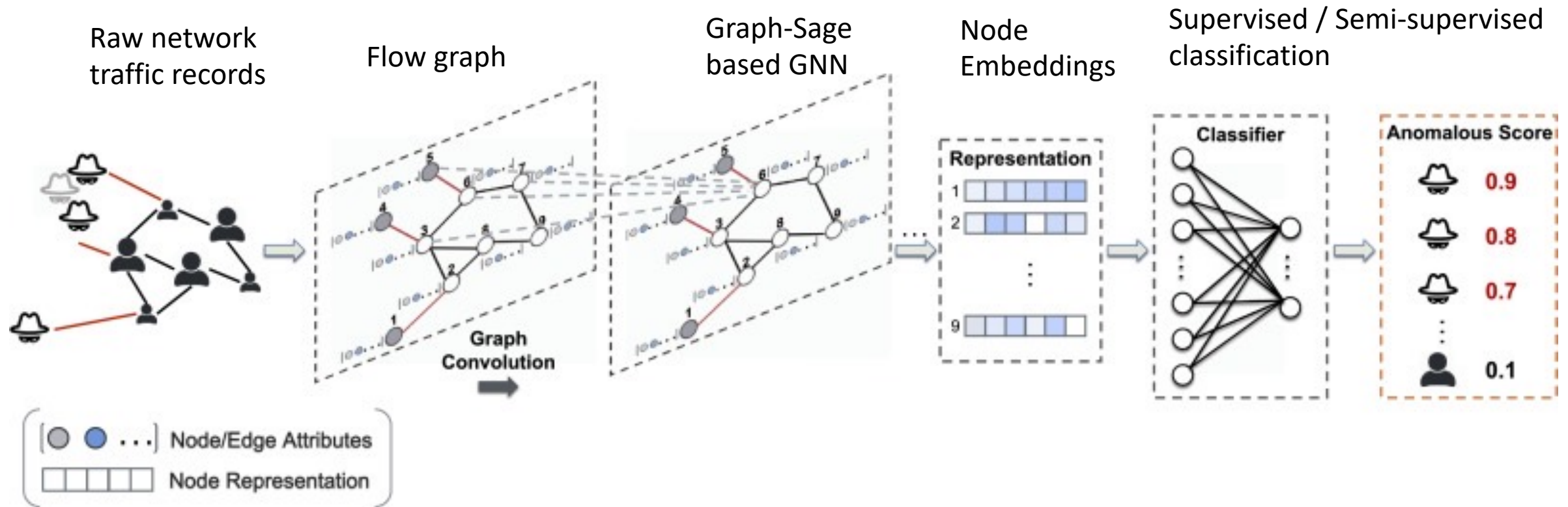
A survey on graph neural networks for intrusion detection systems: Methods, trends and challenges, Computer & Security, Vol.141, June 2024

Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- Target: **Flow Graph** based representation of network traffics

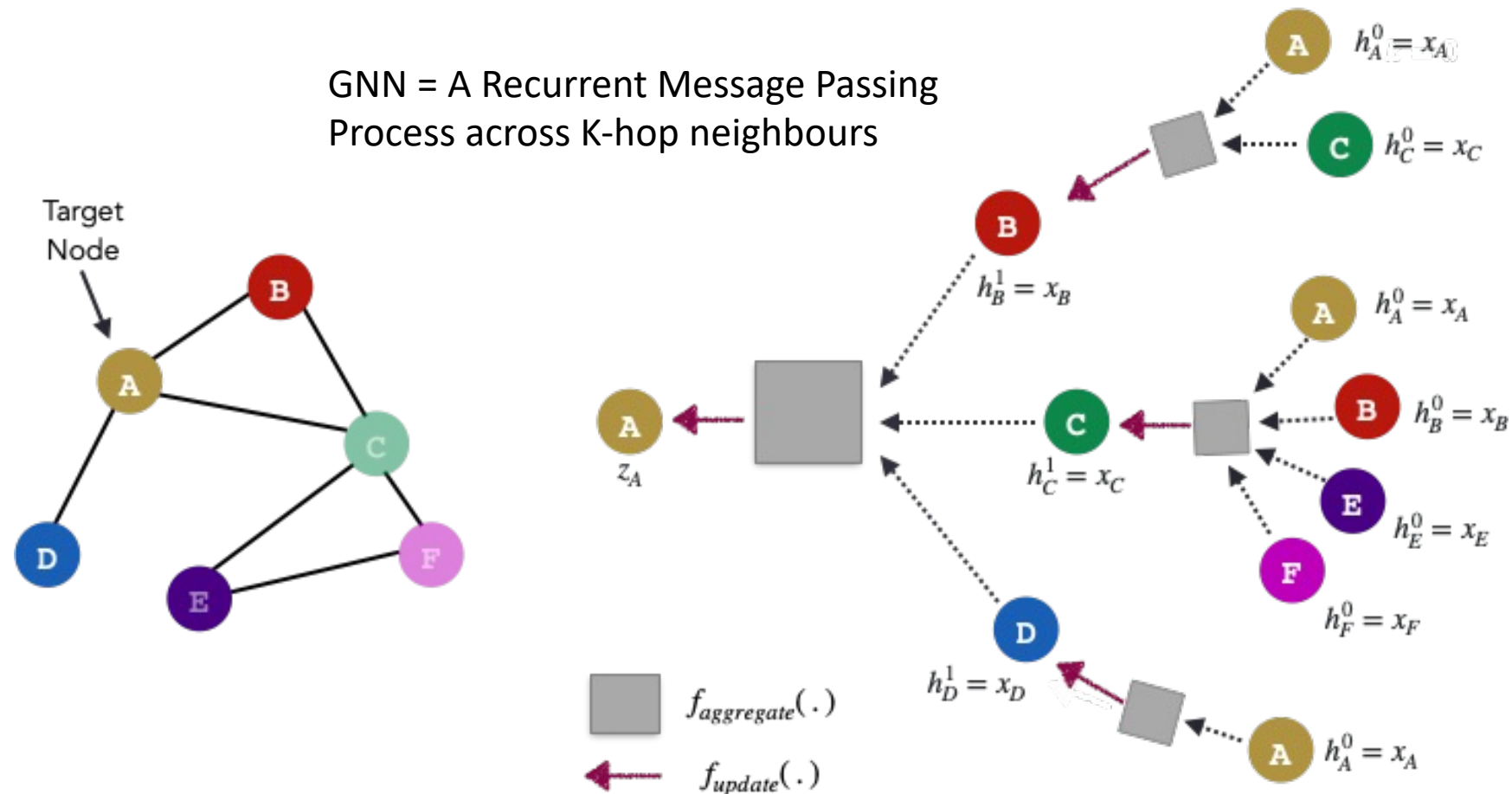


Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection



Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- A few words about Graph-Sage based GNN models



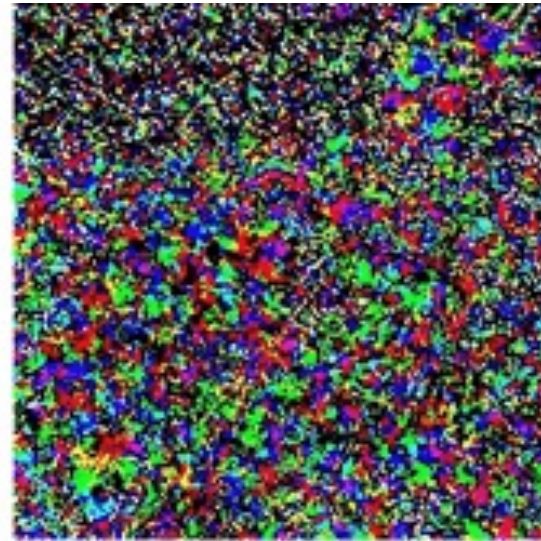
Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- What is an adversarial attack ?



african elephant
probability = 0.83

+ 0.01 x



additive noise

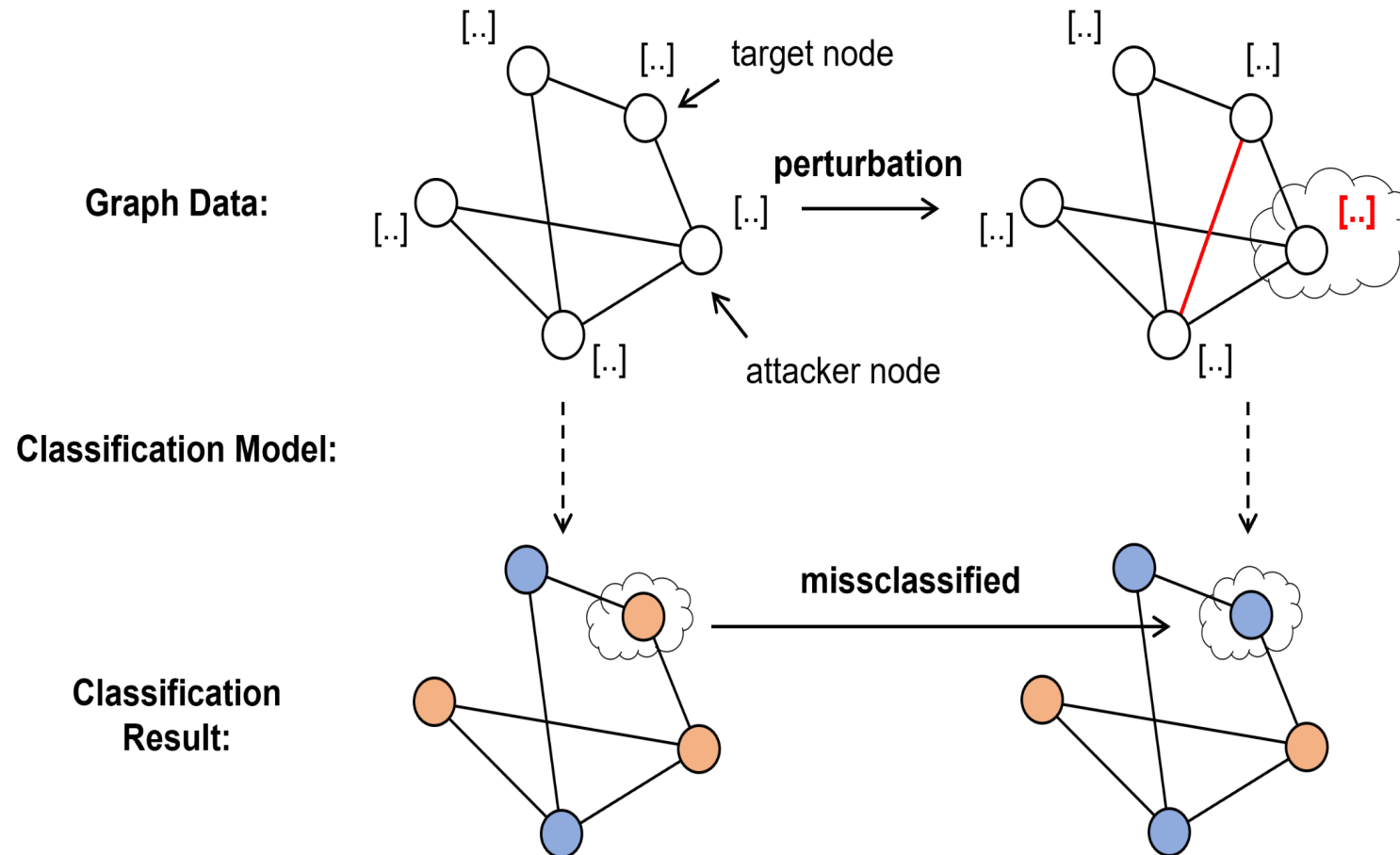
=



sea snake
probability = 0.81

Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- An adversarial attack against Graph Neural Network-based models

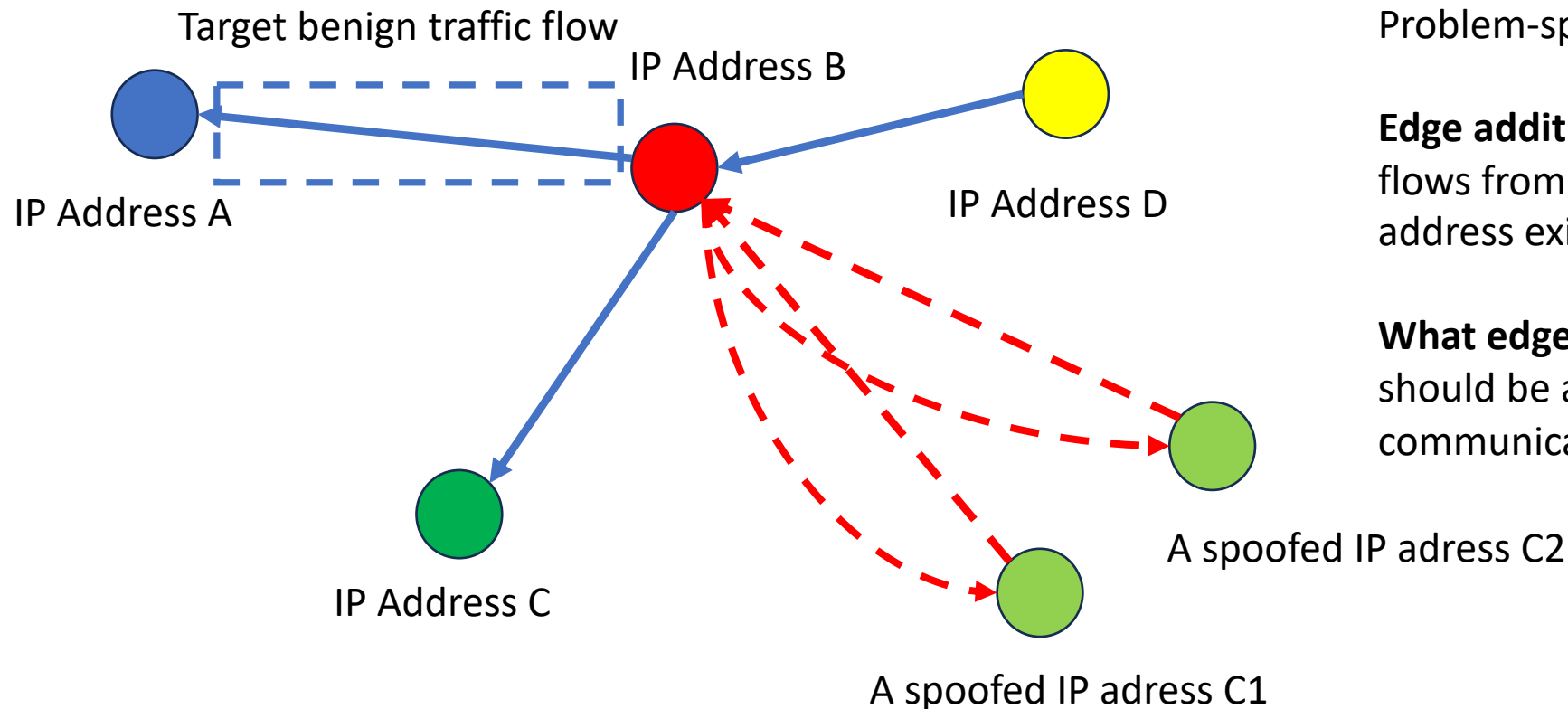


Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- What is the problem space constraint ?
- Any change to the input graph should not break the attack / normal traffic flows
 - An attack is still an attack
- Any change to the input graph should be made consistent with the deployed communication protocols
- In the end, any change to the input graph should be made compatible with the profiles of real traffic flows between IP addresses.

Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- An example of problem-space adversarial attacks



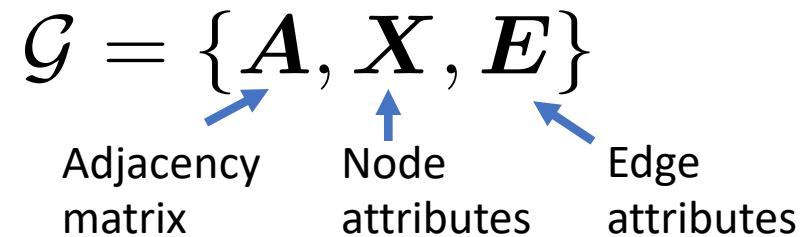
Problem-space constraint:

Edge addition: An attacker can send traffic flows from a spoofed IP address to an IP address existing in the raw traffic records.

What edges to add: The added traffic flows should be added in a way consistent with the communication protocol, such as TCP-IP

Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- Mathematical nature of the edge-modification-based adversarial attacks

$$\mathcal{G} = \{\mathbf{A}, \mathbf{X}, \mathbf{E}\}$$


Adjacency matrix Node attributes Edge attributes

$$\mathbf{A}^* = \arg \max_{\mathbf{A}'} \ell(\mathcal{M}(\mathbf{A}', \mathbf{X}, \mathbf{E}, \mathbf{Y}))$$
$$s.t. \quad \|\mathbf{A} - \mathbf{A}'\|_{L_0} \leq 2b$$

Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- Every GNN is vulnerable to edge modification-based adversarial attacks

Let y be the target class label of the evasion attack. The goal is to make $f_y(\hat{\mathbf{x}})$ deviate from $f_y(\mathbf{x}) = 0$. In other words, we aim at maximizing $f_y(\hat{\mathbf{x}})$, so that $\mathbf{x}$ is modified to get y assigned to it. The evasion attack task can then be formulated as a process of set function optimization, defined as

$$S^* = \arg \max_{|S| \leq K} g(S)$$

$$\text{where } g(S) = \max_{l \in S} f_y(\hat{\mathbf{x}}), \quad l = \text{diff}(\mathbf{b}, \hat{\mathbf{b}})$$

This is in nature a weakly-submodular function maximization problem, which can be solved with Greedy Search with a polynomial complexity.

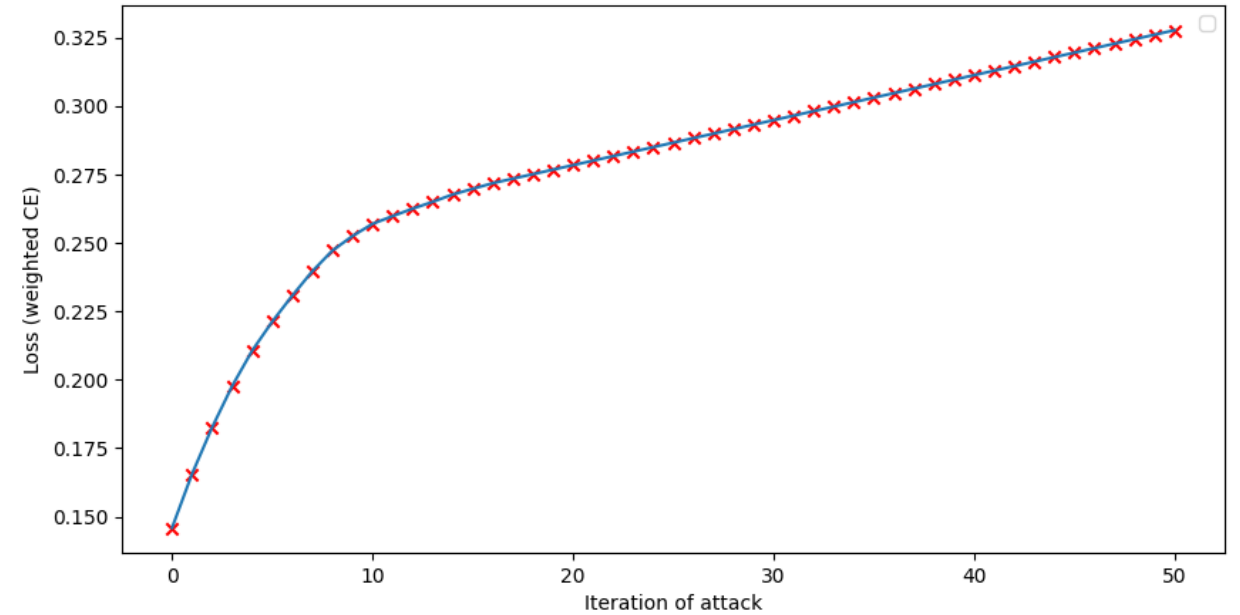
Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

Algorithm 1: Unconstrained Adversarial Attack

Input: Original flow graph $\mathcal{G} = \{\mathbf{A}, \mathbf{X}, \mathbf{E}\}$, attack budget b , target model $\mathcal{M}$, loss function ℓ

Output: Modified flow graph $\mathcal{G}' = \{\mathbf{A}', \mathbf{X}, \mathbf{E}'\}$

```
1 Initialize  $\mathbf{A}' \leftarrow \mathbf{A}$ ,  $\mathbf{E}' \leftarrow \mathbf{E}$ ;  
2 for  $i \leftarrow 1$  to  $b$  do  
3    $max\_loss \leftarrow -\infty$ ;  
4   for each controlled node  $u$  in  $\mathcal{G}$  do  
5     for each possible destination node  $v \neq u$  do  
6       Simulate adding edge  $(u, v)$  to  $\mathbf{A}'$ , with  
7       corresponding edge features  $\mathbf{E}'_{(u,v)}$ ;  
8       Compute loss:  
9        $current\_loss \leftarrow \ell(\mathcal{M}(\mathbf{A}', \mathbf{X}, \mathbf{E}'), \mathbf{Y})$ ;  
10      if  $current\_loss > max\_loss$  then  
11         $max\_loss \leftarrow current\_loss$ ;  
12         $best\_edge \leftarrow (u, v)$ ;  
13   Add  $best\_edge$  to  $\mathbf{A}'$ , update  $\mathbf{E}'$  with its features;  
14 return  $\mathcal{G}' = \{\mathbf{A}', \mathbf{X}, \mathbf{E}'\}$ ;
```



Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- Questions remaining to answer
- We focus on activating an excessively large false alarm rate as the attack goal, but can we also generate evasive samples (increasing the false negative rate ?)
 - Adding benign traffic flows targeting at one IP address in the dataset ?
- Divide the raw network traffics according to their temporal orders.
 - Attackers add adversarial manipulations to historical data, while aiming to deliver adversarial attack impacts over future traffic flows
- The impact of the degree of the target IP address
- Transferability of the attack
 - If an adversarial attack can mislead GNN-based intrusion detection systems, can it also mislead a flow-statistic based intrusion detection model (or not) ?
 - Think about other GNN models

Problem Space-constrained Adversarial Attack against Graph-based Intrusion Detection

- Next step
- Elevating both false alarm and false negative rate via attack
- Where to add edges and what edges to add ?
- Transferable attacks across both the graph-based IDS and statistics-based IDS
- Publication plan:
 - Matthieu Mouzaoui at the second year of this thesis.
 - 2 papers submitted by the end of this year.

Thanks !